



Design and Analysis  
of Algorithms I

# Contraction Algorithm

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## Overview

# Goals for These Lectures

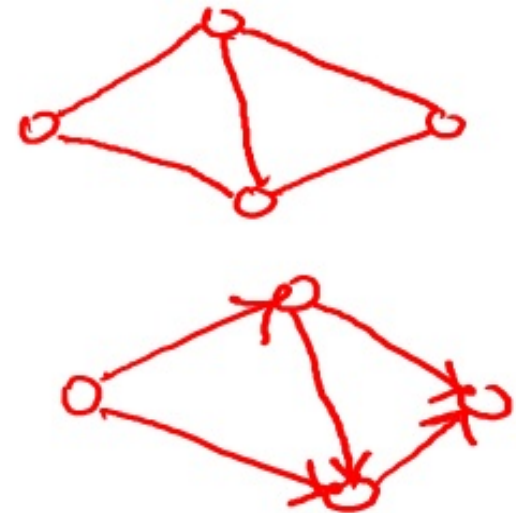
- Further practice with randomized algorithms
  - In a new application domain (graphs)
- Introduction to graphs and graph algorithms

Also: “only” 20 years ago!

# Graphs

## Two ingredients

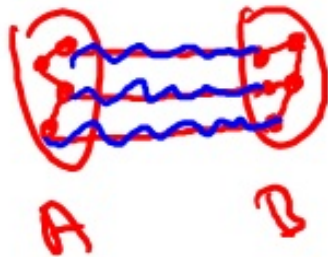
- Vertices aka nodes (V)
- Edges (E) = pairs of vertices
  - can be undirected [unordered pair]
  - or directed [ordered pair] (aka arcs)



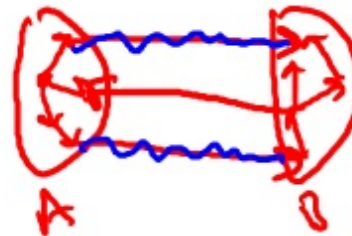
Examples: road networks, the Web, social networks, precedence constraints, etc.

# Cuts of Graphs

Definition: a cut of a graph  $(V, E)$  is a partition of  $V$  into two non-empty sets  $A$  and  $B$ .



[undirected]



[directed]

Definition: the crossing edges of a cut  $(A, B)$  are those with:

- the one endpoint in each of  $(A, B)$  [undirected]
- tail in  $A$ , head in  $B$  [directed]

Roughly how many cuts does a graph with  $n$  vertices have?

☐  $n$


☐  $n^2$

☒  $2^n$

☐  $n^n$

# The Minimum Cut Problem

- INPUT: An undirected graph  $G = (V, E)$ .

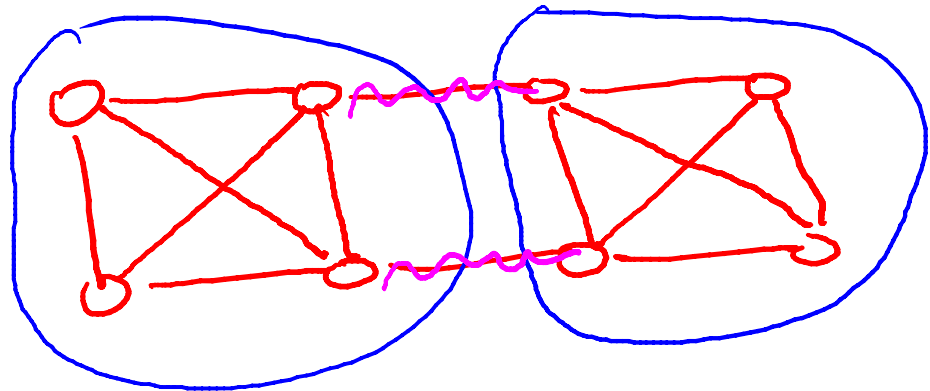
[ Parallel  edges allowed]

[See other video for representation of the input]

- GOAL: Compute a cut with fewest number of crossing edges. (a min cut)

What is the number of edges crossing a minimum cut in the graph shown below?

- ☐ 1
- ☒ 2
- ☐ 3
- ☐ 4



# A Few Applications

- identify network bottlenecks / weaknesses
- community detection in social networks
- image segmentation
  - input = graph of pixels
  - use edge weights
    - [(u,v) has large weight  $\Leftrightarrow$  “expect” u,v to come from some object]

hope: repeated min cuts identifies the primary objects in picture.