



Algorithms: Design
and Analysis, Part II

Huffman Codes

Introduction and Motivation

Binary Codes

Binary code: maps each character of an alphabet Σ to a binary string.

Example: $\Sigma = \text{a-z + various punctuation}$ (size 32 overall, say)

Obvious encoding: use the 32 5-bit binary strings to encode this Σ (a fixed-length code)

Can we do better?: yes, if some characters of Σ are much more frequent than others, using a variable-length code.

Ambiguity

Example: Suppose $S = \{A, B, C, D\}$. Fixed length encoding would be $\{00, 01, 10, 11\}$.

Suppose instead we use the encoding $\{0, 01, 10, 1\}$. What is 001 an encoding of?

- (A) AB → leads to 001
- (B) CD
- (C) AAD → also leads to 001
- (D) not enough info to answer question

Prefix-Free Codes

Problem: with variable-length codes, not clear where one character ends + the next one begins.

Solution: prefix-free codes — make sure that for every pair $i, j \in \Sigma$, neither of the encodings $f(i), f(j)$ is a prefix of the other.

Example: $\{0, 10, 110, 111\}$.

Why useful?: can give shorter encodings with non-uniform character frequencies.

Example

Example:

A	60%	00	0
B	25%	01	10
C	10%	10	110
D	5%	11	111
Σ	frequencies	fixed length	variable-length (prefix-free)

fitted length encoding: 2 bits / character.

Variable length encoding: how many bits needed on average?

- (A) 1.5
- (B) 1.55
- (C) 2
- (D) 2.5

$$\begin{aligned} & .6 \cdot 1 + .25 \cdot 2 \\ & + (.10 + .5) \cdot 3 = \underline{1.55} \end{aligned}$$