

Design and Analysis of Algorithms I

# Graph Algorithms Representing Graphs

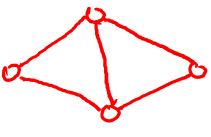
### Graphs

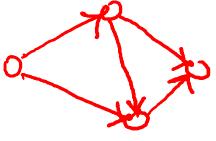
Two ingredients

or directed Cordord pair) (aka arcs)

- vertices alsa nodes (V) - edges (E) = pairs & vertices - Can be undirected (un ordered pair)

Examples! road networks, the Web, social Letworks, precedence constraints, etc.

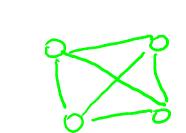




Consider an undirected graph that has n vertices, no parallel edges, and is connected (i.e., "in one piece"). What is the minimum and maximum number of edges that the graph could have, respectively?

$$\bigcirc n-1$$
 and  $n(n-1)/2$ 

- $\bigcirc n-1$  and  $n^2$
- $\bigcirc n$  and  $2^n$
- $\bigcap n$  and  $n^n$



# Sparse vs. Dense Graphs

Let == # of edges.

In most (but not all) applications, mis 52(n) and 0 (n2).

-in a "sparse graph", mis doser to O(n2)

# The Adjacency Matrix

Represent C by a non 0-1 matrix A, where

A:=1 => 6 hes an i-jedge (D-C)

How much space does an adjacency matrix require, as a function of the number n of vertices and the number m of edges?

- $\bigcirc \theta(n)$
- $\bigcirc \theta(m)$
- $\bigcirc \theta(m+n)$
- $\bigcirc \theta(n^2)$

### **Adjacency Lists**

```
Ingredients

- array (or list) of vertices

- array (or list) of edges

- each edge points to its endpoints

- each vertex points to edges incident on it
```

How much space does an adjacency list representation require, as a function of the number n of vertices and the number m of edges?

- $\bigcirc \theta(n)$
- $\bigcirc \theta(m)$
- $\bigcirc \theta(m+n)$ 
  - $\bigcirc \theta(n^2)$

Adjacency Lists

Space Ingredients (n) O - array (or list) of vertices (m) ( - array (or list) of edges - each edge points to its endpoints O CWJ - each vertex points to edges incident on it 9 (m) (m+n) Question: which is better? [or O(wax {w "3)] Answer: depends on graph density and operations needed. This course: Fo cus adjacency lists.

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