



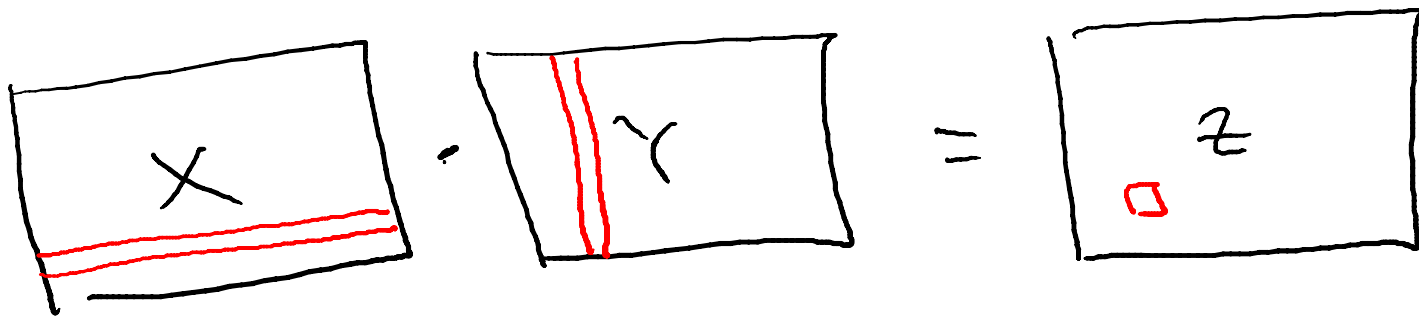
Design and Analysis  
of Algorithms I

# Divide and Conquer

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# Matrix Multiplication

# Matrix Multiplication



(all  $n \times n$  matrices)

where  $z_{ij} = (\text{ith row of } X) \cdot (\text{jth column of } Y)$

$$= \sum_{k=1}^n x_{ik} \cdot y_{kj}$$

Note: input size  
 $\in \mathcal{O}(n^2)$

## Example (n=2)

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} e & f \\ g & h \end{pmatrix} = \begin{pmatrix} ae+bg & af+bh \\ ce+dg & cf+dh \end{pmatrix}$$

$$z_{ij} = \frac{\sum_{k=1}^n x_{ik} y_{kj}}{\underbrace{\hspace{10em}}} \rightarrow O(n)$$

What is the asymptotic running time of the straightforward iterative algorithm for matrix multiplication?

☐  $\theta(n \log n)$

☐  $\theta(n^2)$

 ☒  $\theta(n^3)$

☐  $\theta(n^4)$

# The Divide and Conquer Paradigm

- ① DIVIDE into smaller subproblems.
- ② CONQUER subproblems recursively.
- ③ COMBINE solutions of subproblems into one for the original problem.

# Applying Divide and Conquer

Idea:  
with  $X = \begin{pmatrix} A & B \\ C & D \end{pmatrix}$  and  $Y = \begin{pmatrix} E & F \\ G & H \end{pmatrix}$

[where A through H are all  $\frac{n}{2} \times \frac{n}{2}$  matrices]

Then: (you check)

$$X \cdot Y = \begin{pmatrix} AE + BG & AF + BH \\ CE + DG & CF + DH \end{pmatrix}$$

# Recursive Algorithm #1

$$X \cdot Y = \begin{pmatrix} AE + BG & AF + BH \\ CE + DG & CF + DH \end{pmatrix}$$

Step 1: recursively compute the 8 necessary products.

Step 2: do the necessary additions  
( $\Theta(n^2)$  time)

Fact: run time is  $\Theta(n^3)$ . [follows from master method]

# Strassen's Algorithm (1969)

Step 1: recursively compute only 7  
(cleverly chosen) products

Step 2: do the necessary (clever) additions  
+ subtractions (still  $\Theta(n^2)$  time)

Fact: better than cubic time!  
(see Master Method lecture)



$$X = \begin{pmatrix} A & B \\ C & D \end{pmatrix}$$

## The Details

$$Y = \begin{pmatrix} E & F \\ G & H \end{pmatrix}$$

The Seven Products :  $P_1 = A(F-H)$ ,  $P_2 = (A+B)H$ ,  
 $P_3 = (C+D)E$ ,  $P_4 = D(G-E)$ ,  $P_5 = (A+D)(E+H)$ ,  
 $P_6 = (B-D)(G+H)$ ,  $P_7 = (A-C)(E+F)$

Claim :  $X \cdot Y = \begin{pmatrix} AE+BG & AF+BH \\ CE+DG & CF+DH \end{pmatrix} = \begin{pmatrix} P_5+P_4-P_2+P_6 & P_1+P_2 \\ P_3+P_4 & P_1+P_5-P_3-P_7 \end{pmatrix}$

Proof :  $AE + \cancel{AH} + \cancel{DE} + \cancel{DH} + \cancel{DG} - \cancel{DE} - \cancel{AH} - \cancel{BH}$   
 $+ \cancel{BG} + \cancel{BH} - \cancel{DG} - \cancel{DH} = AE + BG$

Question: where did this come from?

(remains open!)