Design and Analysis of Algorithms I

Introduction

Merge Sort (Pseudocode)

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Merge Sort: Pseudocode

- recursively sort 1st half of input array
- " " 2nd " " " "
- merge two sorted sublists into one

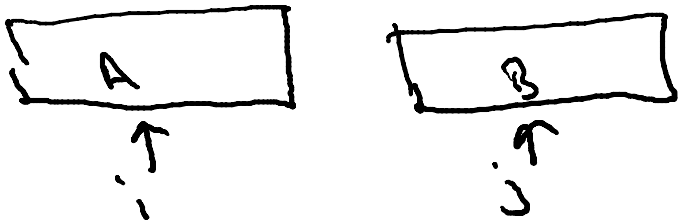
[ignores base cases]

Pseudocode for Merge:

C = output array [length = n]

A = 1st sorted array [n/2]

B = 2nd " " [n/2]



i = 1

j = 1

```
for k = 1 to n
    if A(i) < B(j)
        C(k) = A(i)
        i++
    else [B(j) < A(i)]
        C(k) = B(j)
        j++
end [ignore end cases]
```

Pseudocode for Merge:

```
C = output [length = n]
A = 1st sorted array [n/2]
B = 2nd sorted array [n/2]
i = 1
j = 1
```

```
for k = 1 to n
        if A(i) < B(j)
                C(k) = A(i)
                i++
        else [B(j) < A(i)]
                C(k) = B(j)
                j++
end
```

(ignores end cases)

Merge Sort Running Time?

Key Question: running time of MergeSort on array of n numbers?

[running time $\approx$ # of lines of code executed]

Pseudocode for Merge:

C = output [length = n]

A = 1st sorted array [n/2]

B = 2nd sorted array [n/2]

i = 1

j = 1

} 2 operations

```
for k = 1 to n
        if A(i) < B(j)
                C(k) = A(i)
                i++
        else [B(j) < A(i)]
                C(k) = B(j)
                j++
end
```

(ignores end cases)

Running Time of Merge

Upshot: running time of Merge on array of m numbers is $\leq 4m + 2$

$\leq 6m$ (since $m \geq 1$)

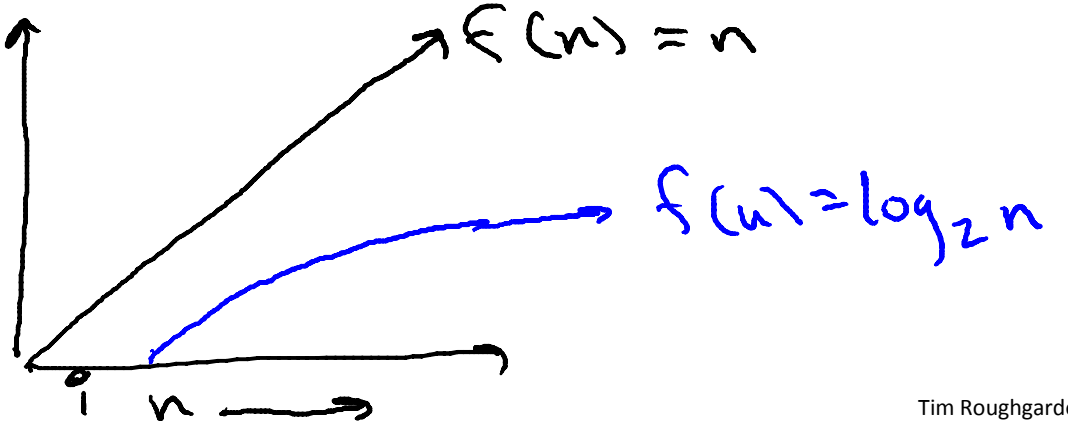
Running Time of Merge Sort

Claim: MergeSort requires

$\leq 6n \log_2 n + 6n$ operations

to sort n numbers.

Recall: $\log_2 n$ = # of times you divide by 2 until you get down to 1.



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