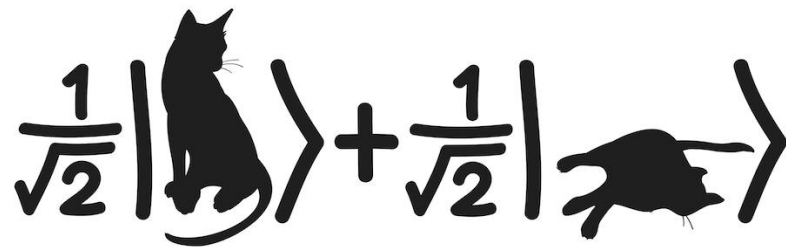


Quantum Mechanics & Quantum Computation

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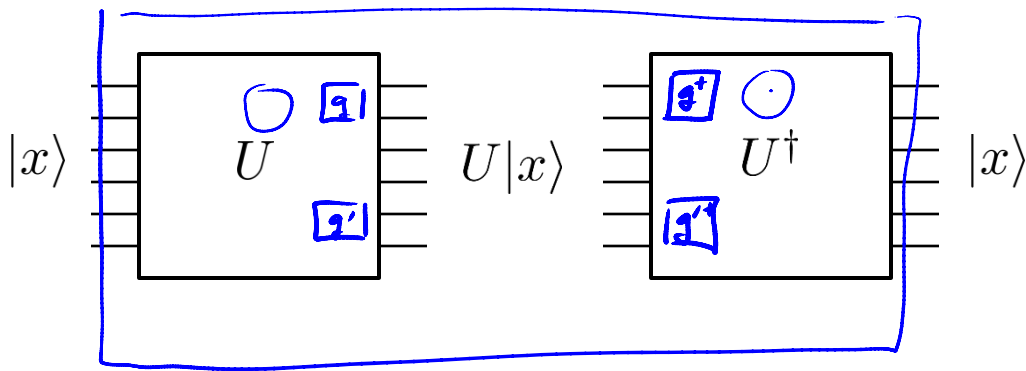


Lecture 11: Quantum Circuits

Reversible Computation

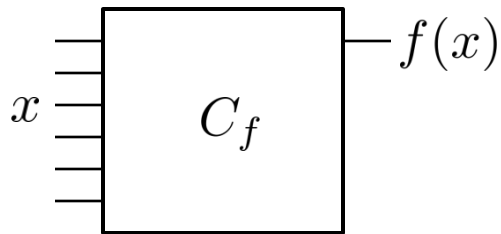
Reversible computation

- Quantum computers are reversible.
- Why?



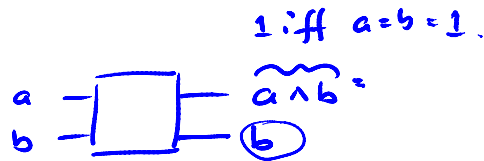
$$UU^\dagger = U^\dagger U = I.$$

Implementing classical circuits

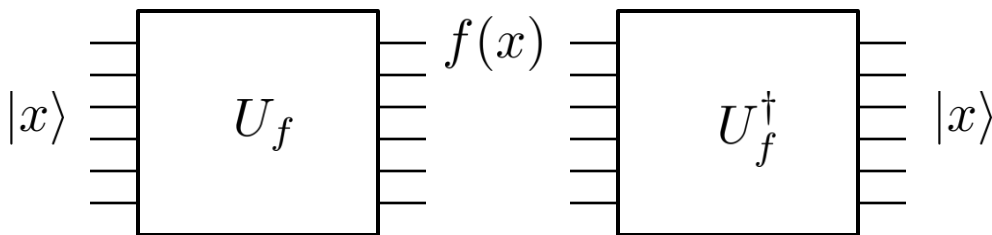


Classical circuit for computing a boolean function

$$f : \{0, 1\}^n \rightarrow \{0, 1\}$$



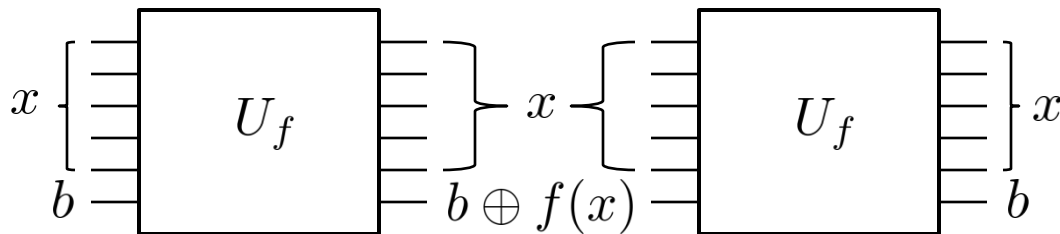
Imagine a quantum version of C_f :



Have to be reversible.

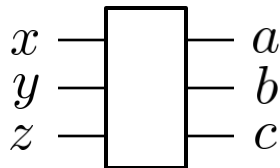
But classical gates throw away information!

Classical reversible computation



$b=0$

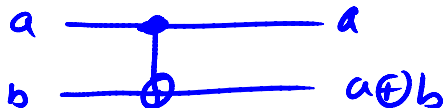
Consider C-SWAP gate:



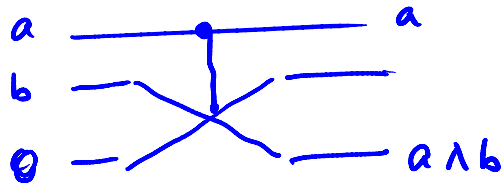
$$\begin{aligned} x = 0 &\Rightarrow a = x, b = y, c = z \\ x = 1 &\Rightarrow a = x, b = z, c = y \end{aligned}$$

NOT $|0\rangle \rightarrow \boxed{x} \rightarrow |1\rangle$

CNOT

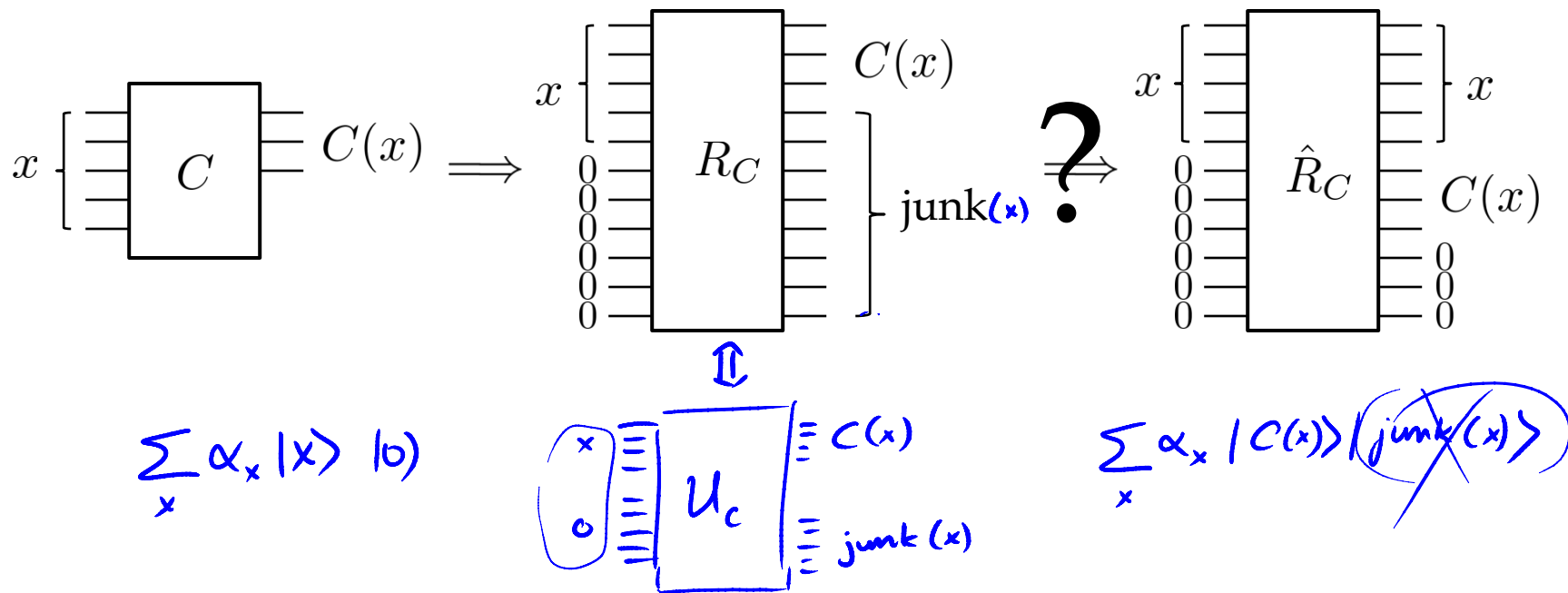


C-SWAP



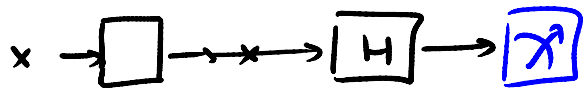
if $a=0$	0
if $a=1$	b

Classical reversible computation



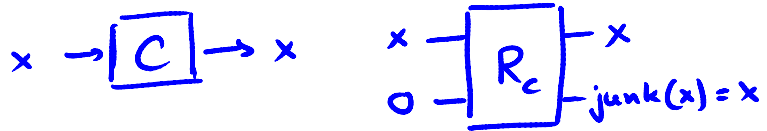
Why remove junk?

Prevents interference.



$$\sum \alpha_x |x\rangle \rightarrow \sum \alpha_x |x\rangle$$

$$\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle \rightarrow |0\rangle$$



$$\begin{pmatrix} \frac{1}{\sqrt{2}} |0\rangle \\ \frac{1}{\sqrt{2}} |1\rangle \end{pmatrix} \xrightarrow{H} \begin{pmatrix} \frac{1}{2} |0\rangle + \frac{1}{2} |1\rangle \\ \frac{1}{2} |0\rangle - \frac{1}{2} |1\rangle \end{pmatrix}$$

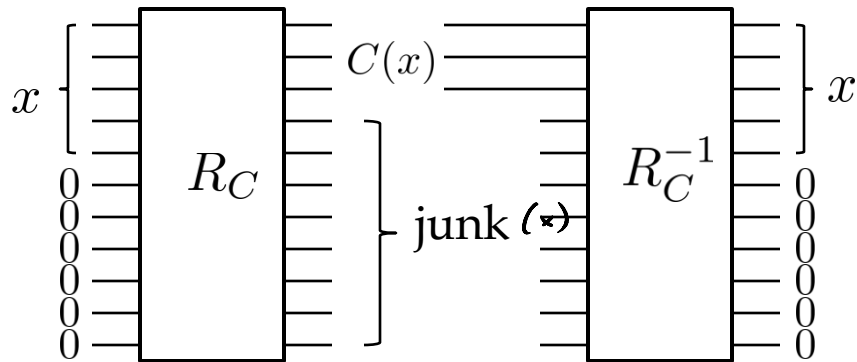
$\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$
 $0 -$
 $\frac{1}{2}|00\rangle + \frac{1}{2}|10\rangle + \frac{1}{2}|01\rangle - \frac{1}{2}|11\rangle$
 $\frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle$

Why remove junk?

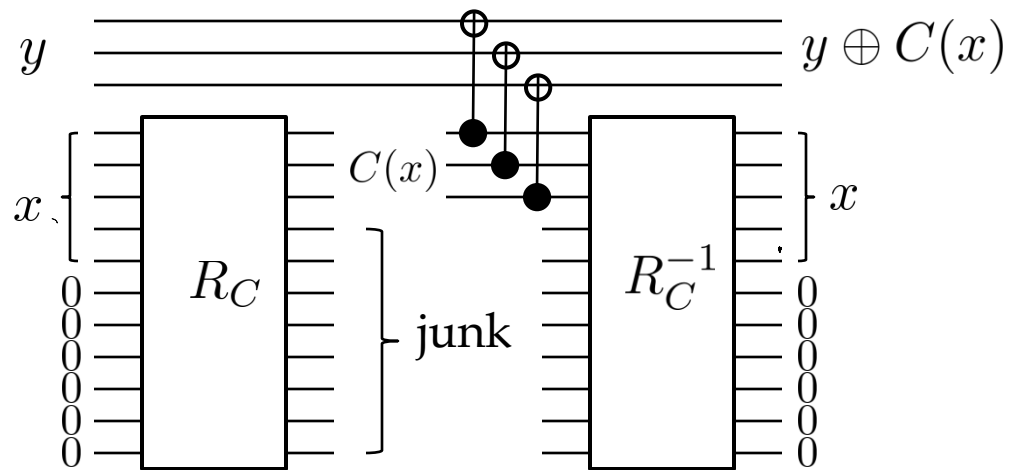
- Can't we just throw away the junk qubits?

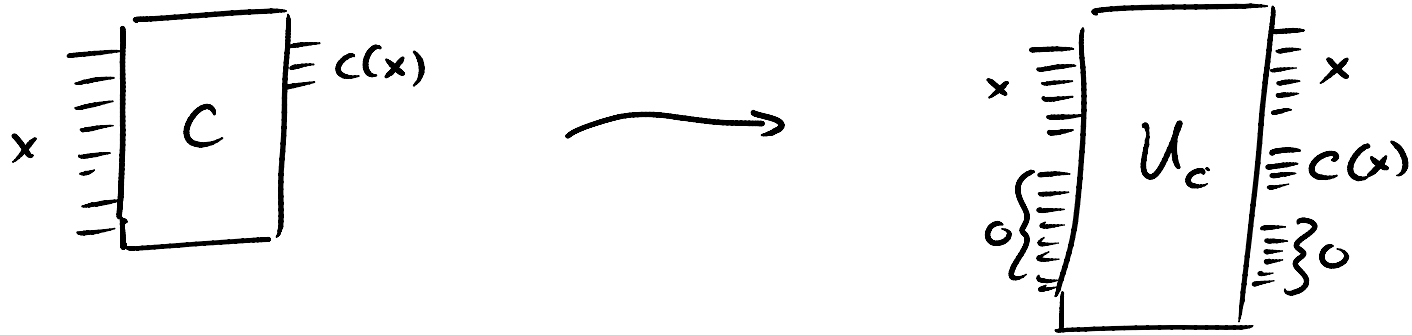
$$\begin{array}{c} \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle \\ |0\rangle \end{array} \begin{array}{|c} \hline \text{ } \\ \hline \end{array} \begin{array}{c} \frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle \end{array}$$

Classical reversible computation



Classical reversible computation





$$\sum_x \alpha_x |x\rangle |0\dots 0\rangle \xrightarrow{U_C} \sum_x \alpha_x |x\rangle |c(x)\rangle |0\dots 0\rangle$$

$$|xy\rangle = |x\rangle |y\rangle = |x\rangle \otimes |y\rangle$$